

# Near-optimal Batch Mode Active Learning and Stochastic Optimization

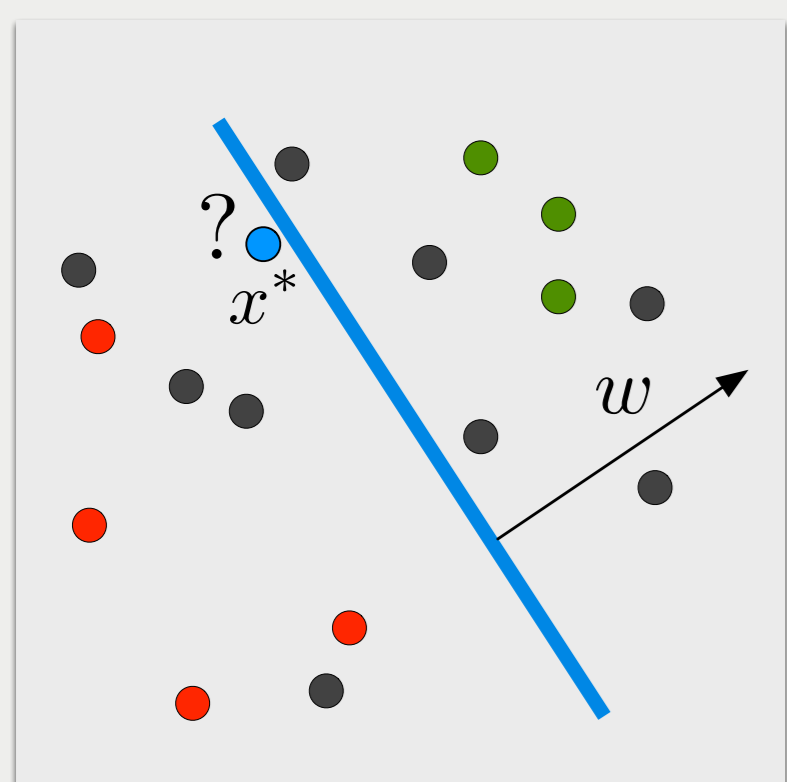
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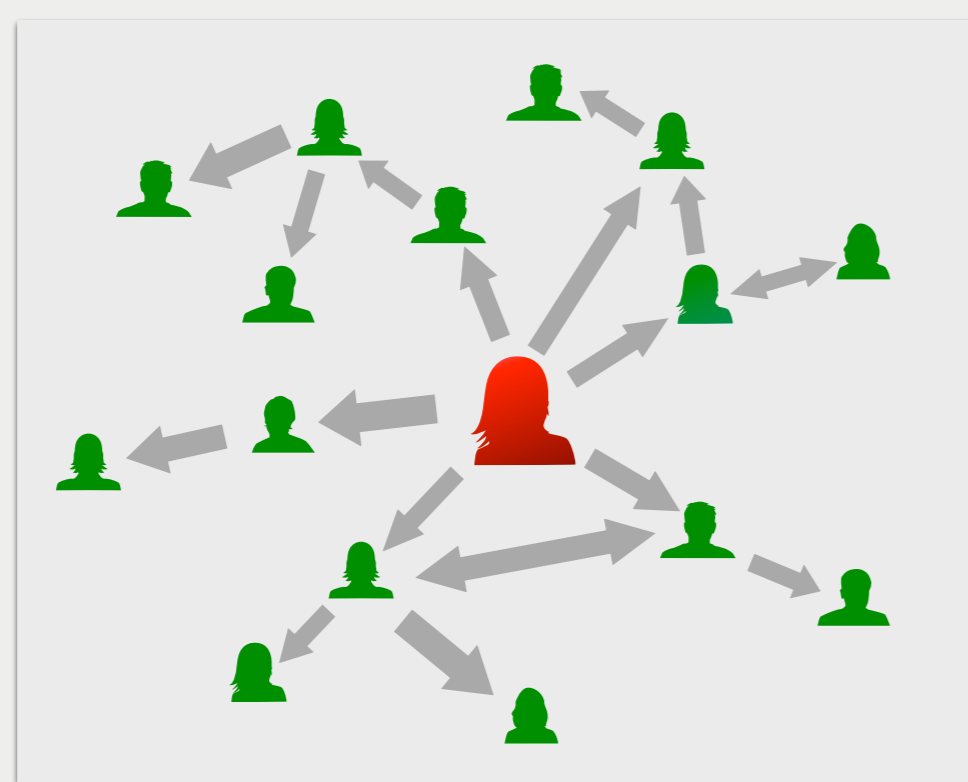
## Contributions

- A general approach for information parallel learning and decision making
- Strong performance guarantees for a simple BATCHGREEDY algorithm
- Practical for batch-mode active learning and influence maximization
- Demonstrate the effectiveness of the algorithms for both applications

## Adaptive Stochastic Minimum Cost Cover (Golovin & Krause '11)



[A]ctive learning



[I]nfluence maximization

**A**:  $f$ : # of hypotheses eliminated  
 $\mathbf{y}_V$ : labels of data points  
 $P$ : hypotheses prior

**I**:  $f$ : # of nodes influenced  
 $\mathbf{y}_V$ : Influences of nodes  
 $P$ : Joint distribution of  $\mathbf{y}_V$

$$\min_{\pi \in \Pi} \text{cost}(\pi) \text{ s.t. } f(\mathcal{S}(\pi, \mathbf{y}_V)) \geq Q \text{ for all } \mathbf{y}_V \text{ with } P(\mathbf{y}_V) > 0$$

## The Min-Cost Cover subjective

**Normalized** We derive no utility from knowing nothing:  $f(\emptyset) = 0$

**Monotonic** Adding labels never hurts:

$$\mathcal{S} \subseteq \mathcal{S}' \subseteq \mathcal{V} \times \mathcal{O} \Rightarrow f(\mathcal{S}) \leq f(\mathcal{S}')$$

**Submodular** Adding a label helps more if we have observed less labels:

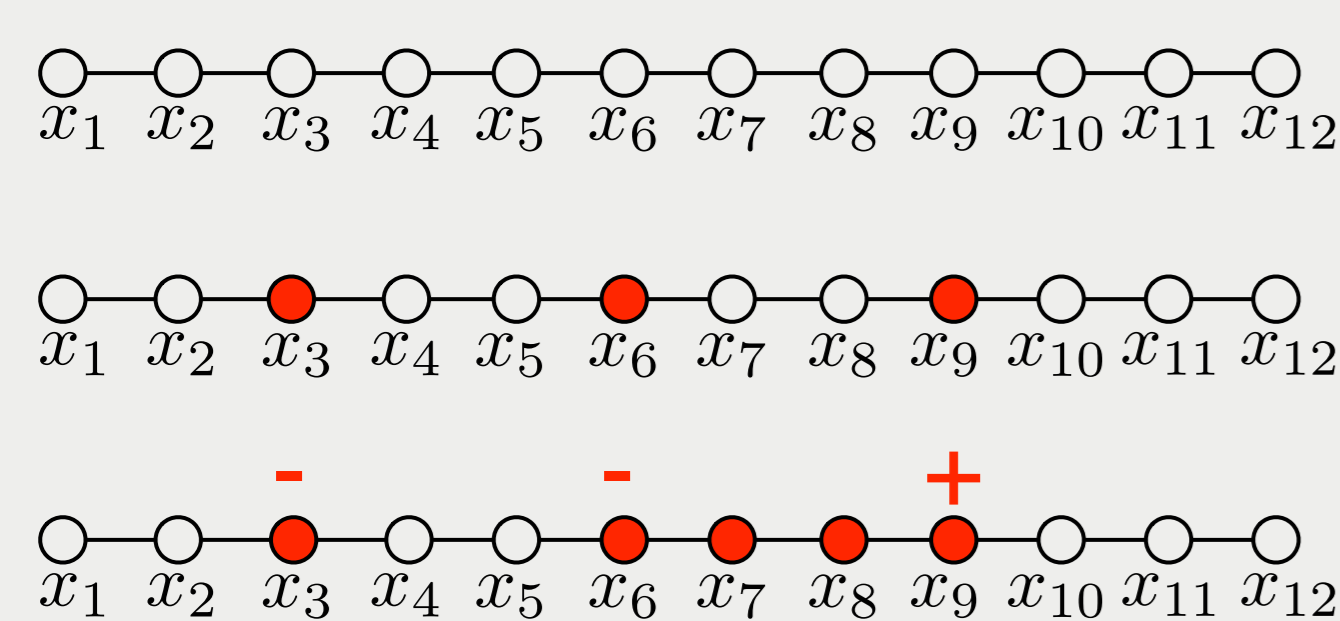
$$\left. \begin{array}{l} \mathcal{S} \subseteq \mathcal{S}' \subseteq \mathcal{V} \times \mathcal{O} \\ (j, y) \in \mathcal{V} \times \mathcal{O} \setminus \mathcal{S}' \end{array} \right\} \Rightarrow f(\mathcal{S} \cup \{(j, y)\}) - f(\mathcal{S}) \geq f(\mathcal{S}' \cup \{(j, y)\}) - f(\mathcal{S}')$$

**Adaptive submodular** The gain of an item, in expectation over its unknown label, can never increase as we gather more information:

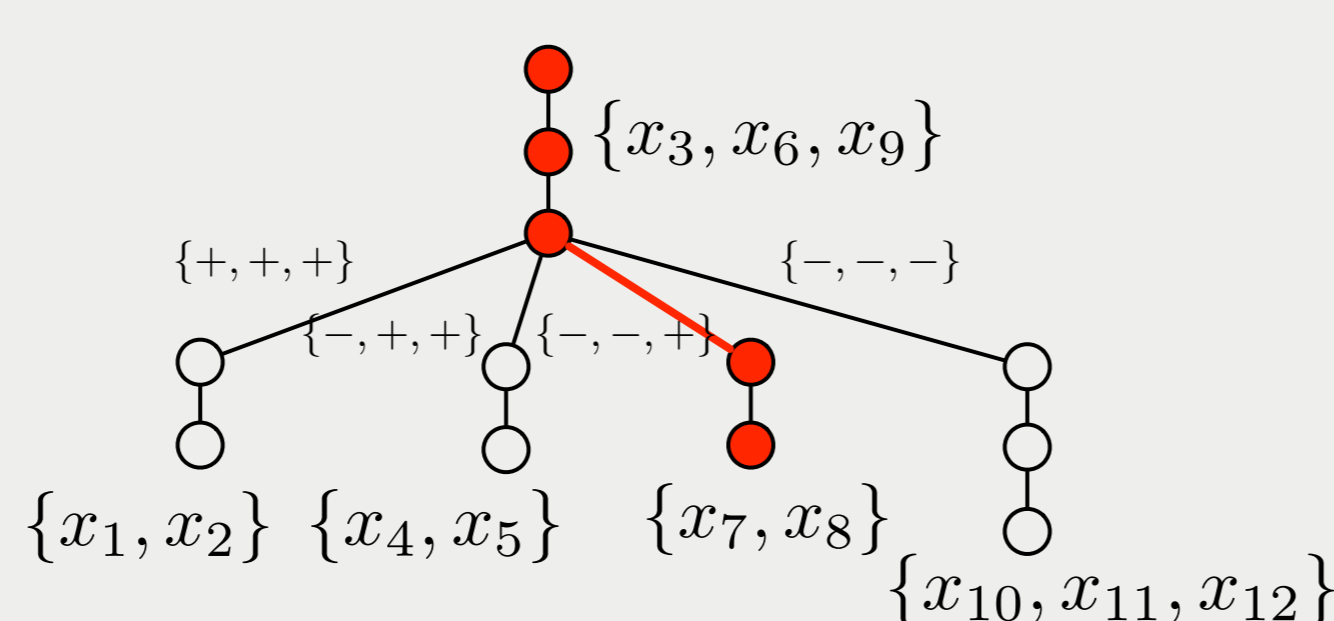
$$\mathcal{S} \subseteq \mathcal{S}' \subseteq \mathcal{V} \times \mathcal{O}, P(\mathcal{S}') > 0 \Rightarrow \Delta_f(j | \mathcal{S}) \geq \Delta_f(j | \mathcal{S}')$$

$$\text{where } \Delta_f(j | \mathcal{S}) = \sum_y P(Y_j = y | \mathcal{S}) [f(\mathcal{S} \cup \{(j, y)\}) - f(\mathcal{S})]$$

## Illustration of Batch Mode Active Learning (1-d, binary threshold)



(top) unlabeled data; (middle) 1st selected batch; (bottom) received labels & 2nd selected batch.



The decision tree representing the BATCHGREEDY policy

## Near-optimal Batch Selection Algorithm (BatchGreedy)

Conditional marginal benefit of item  $s$ :

$$\Delta_f(s | \mathcal{A}, \mathbf{y}_B) = \mathbb{E}_{\mathbf{y}_V} [f(\mathbf{y}_{\mathcal{S} \cup \{s\}} \cup \mathcal{A} \cup \mathbf{y}_B) - f(\mathbf{y}_{\mathcal{A} \cup \mathbf{y}_B}) | \mathbf{y}_B]$$

The BATCHGREEDY policy will greedily select the  $i$ -th element in the  $j$ -th batch

$$s_{i,j} = \arg \max_{s \in \mathcal{V}} \Delta_f(s | \{s_{1,j}, \dots, s_{i-1,j}\}, \mathbf{y}_B)$$

where  $\mathbf{y}_B$  is the set of observations (labeled examples) from batches up to  $j - 1$

## Comparing BatchGreedy with the Optimal Batch Policy

Let  $\text{OPT}_{ac,k}$  be the expected cost and  $\text{OPT}_{wc,k}$  be the worst-case cost of an optimal policy selecting batches of size  $k$ . Further let  $\delta = \min_{\mathbf{y}_V \in \text{supp}(P)} P(\mathbf{y}_V)$ . Then for the cost of the policy  $\pi_G$  implementing BATCHGREEDY it holds that

$$\text{cost}(\pi_G) \leq \text{OPT}_{ac,k} \left( \frac{e}{e-1} \right) (\ln Q + 1), \text{ and}$$

$$\text{cost}_{wc}(\pi_G) \leq \text{OPT}_{wc,k} \left( \frac{e}{e-1} \right) \left( \ln \frac{Q}{\delta} + 1 \right).$$

## Comparing BatchGreedy with the Optimal Sequential Policy

Fix  $0 < \beta < 1$ . Let  $\text{OPT}_{wc}$  be the worst-case cost of an optimal *sequential* policy  $\pi^*$ , constrained to picking a number of items which is a multiple of  $k$ . Further suppose that the variables  $Y_1, \dots, Y_n$  are **independent**. Then for the cost of the policy  $\pi_G$  implementing BATCHGREEDY, run until it achieves  $f(\pi_G) \geq Q - \beta$  it holds that

$$\text{cost}_{wc}(\pi_G) \leq \text{OPT}_{wc} \left( \frac{e}{e-1} \right)^2 \left( \ln \frac{Q}{\beta} + 1 \right).$$

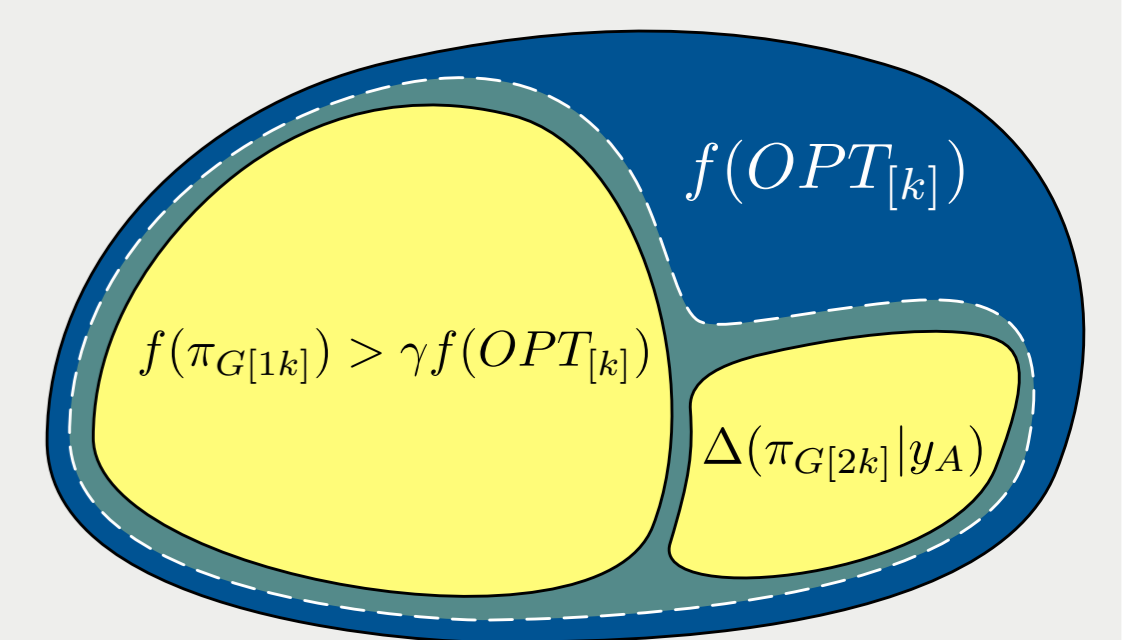
Moreover, it holds that  $P(f(\mathcal{S}(\pi_G, \mathbf{y}_V)) \geq Q) \geq 1 - \beta$ .

## BatchGreedy VS. the Optimal Sequential Policy

### Single batch setting

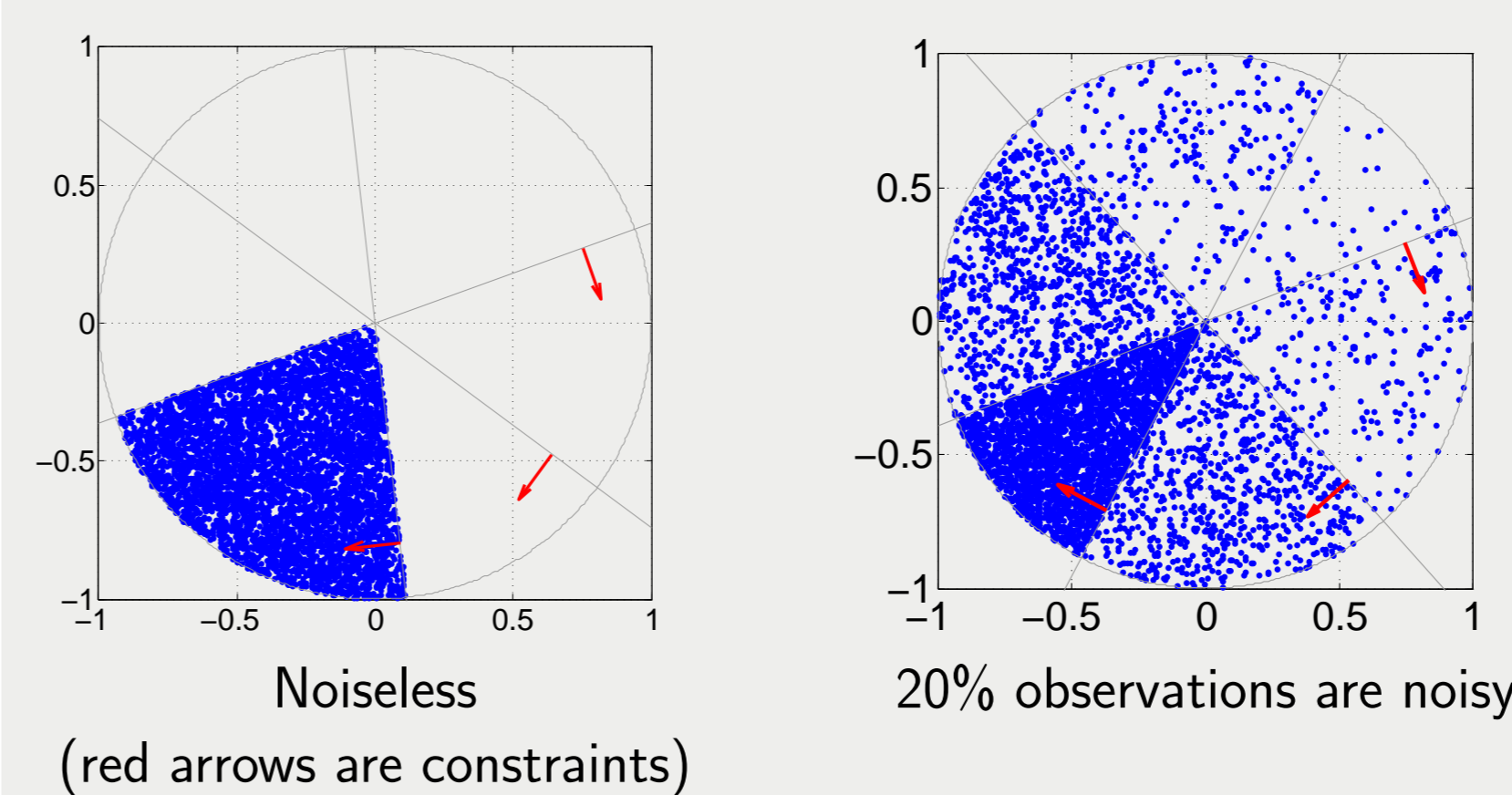
- The expected **cost** incurred by sequential policy to achieve  $Q$  can be *arbitrarily lower* than the cost of single-batch setting (Goemans and Vondrak '06).
- The **utility** of sequential algorithm (with cardinality cstr.) *cannot be much higher* than that of the best non-adaptive algorithm (Asadpour et al, '08).

**Key insight** We can string multiple batches together in an adaptive manner, in a way that the overall **cost** of achieving quota  $Q$  is not much higher than that of the fully sequential policy.



## An Approximate Implementation of BatchGreedy for BMAL

Running time of BATCHGREEDY depends polynomially on  $|\mathcal{H}|$ .



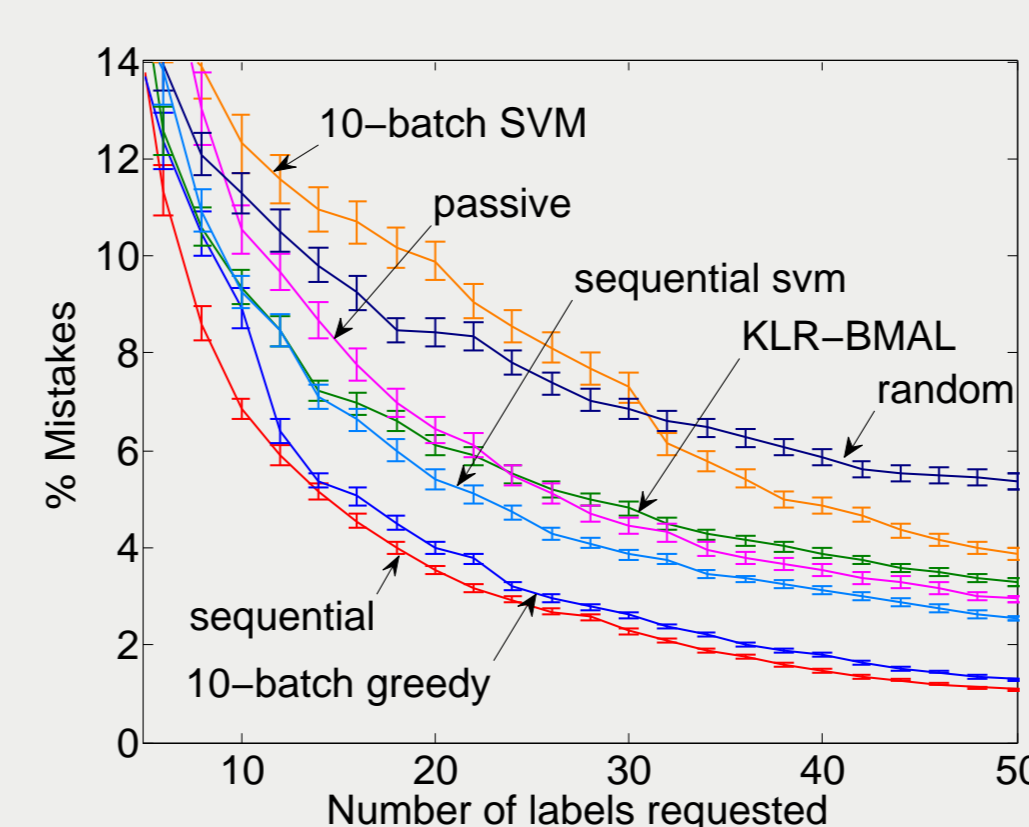
**MCMC sampler** Hit-and-run hypothesis sampler for linear separators and noisy observations

**Assumption** Hypotheses that violate more constraints induce lower confidence

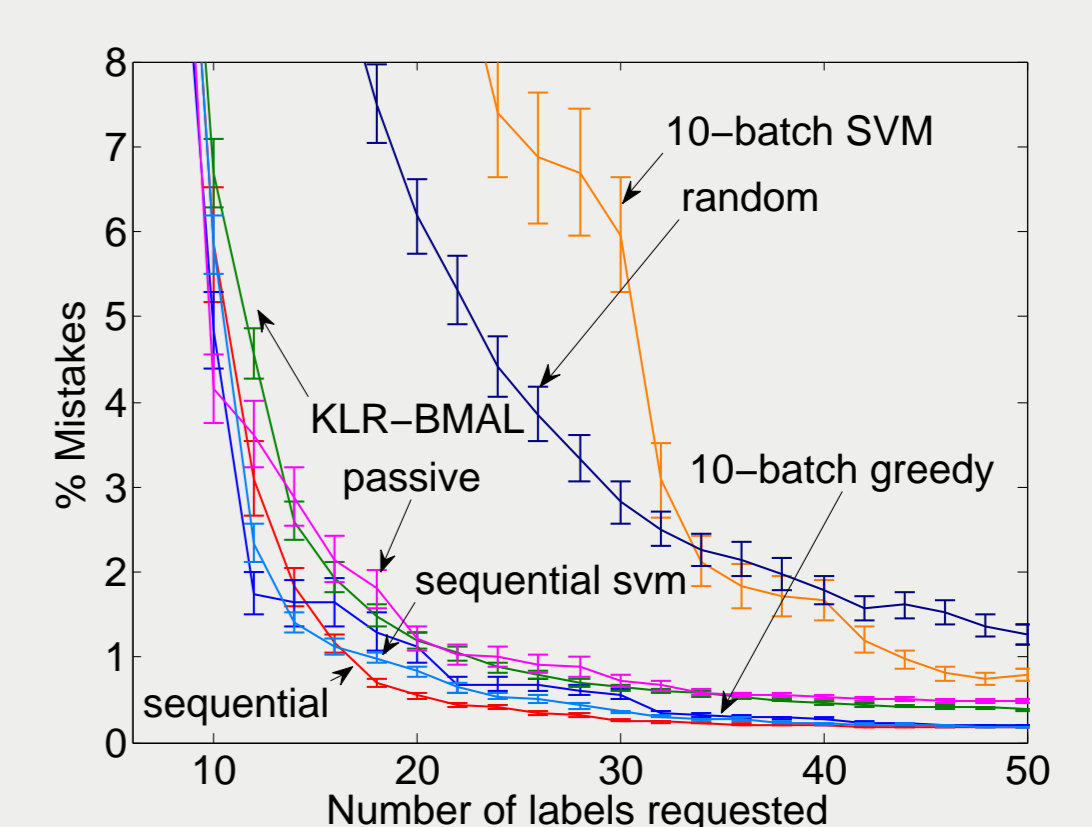
$$s \leftarrow \arg \min_{s'} \sum_{\ell=1}^N \left[ \hat{P}(\mathcal{H}(\{(x, h_{\ell}(x)) : x \in \mathcal{A} \cup \{s'\}\})) \right]$$

**Between batches** Update  $\hat{P}(\mathcal{H}) = \frac{1}{N} \sum_{\ell=1}^N \delta_{h_{\ell}}$  based on current observations.

## Pool-based Batch Mode Active Learning



WDBC

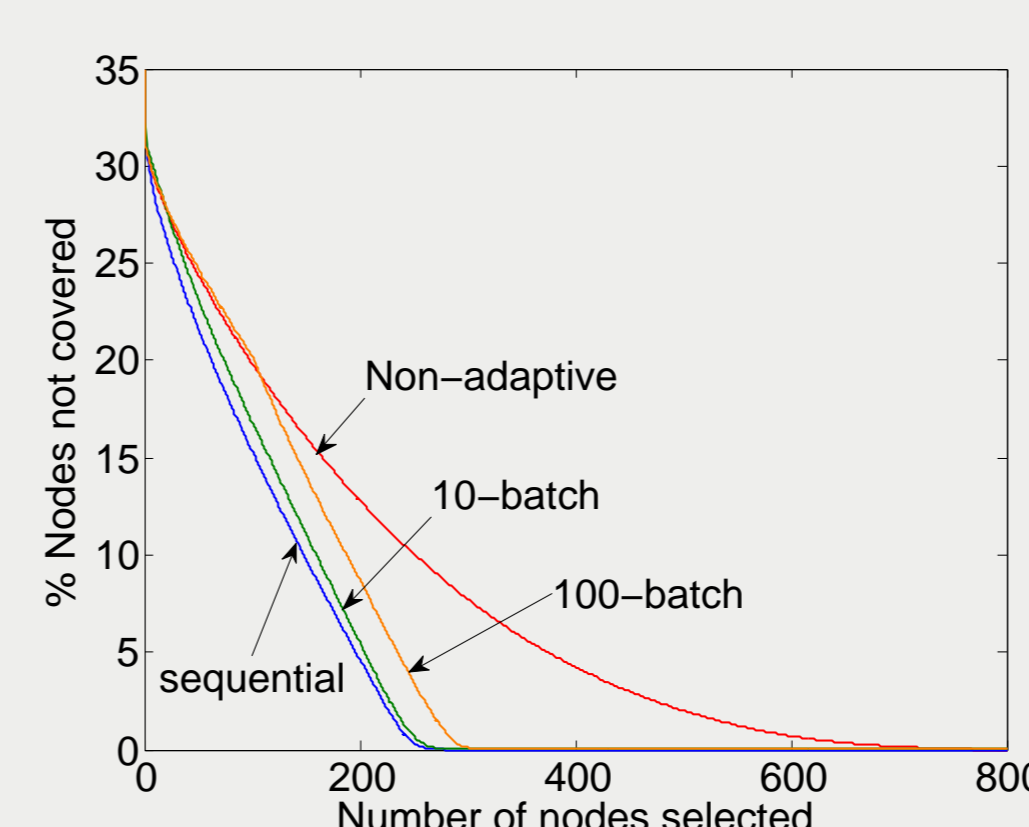


MNIST

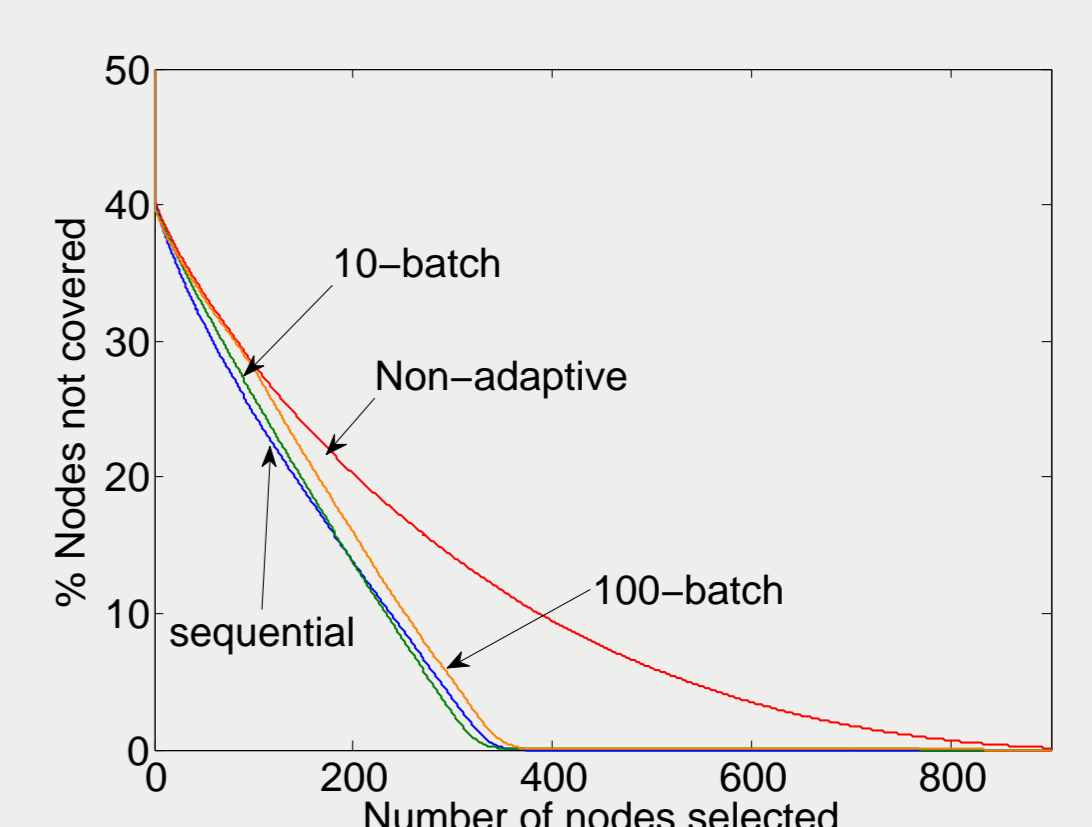
## Applications

- Crowdsourcing annotation on Amazon Mechanical Turk
- High-throughput experimental design

## Multi-stage Influence Maximization in Social Networks



Epinions



Slashdot

## Applications

- Multi-stage marketing campaign
- Resource allocation in computational sustainability
- Vaccination problems in epidemiology